

## Dynamic Programming

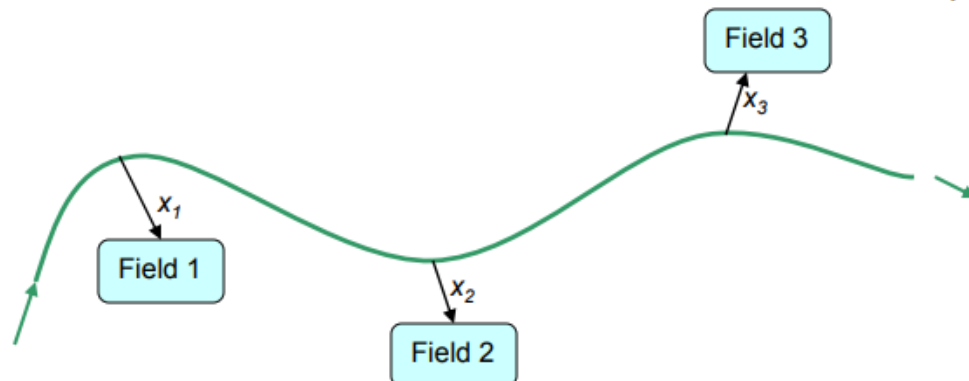
**Dynamic programming** provides an extremely powerful approach for solving the optimization problems that occur in the operation and planning of **water resource** systems. Dynamic programming provides an extremely powerful approach for solving the optimization problems that occur in the operation and planning of water resource systems. However, its applicability has been somewhat limited because of the large computational requirements of the standard algorithm.

**Dynamic Programming** : A Sequential or multistage decision making process .

Water Allocation problem is solved as a sequential process using **dynamic** approach.

programming Objectives :-(1) To discuss the Water Allocation Problem (2) To explain and develop **recursive equations** for backward approach (3) To explain and develop recursive equations for forward approach

Consider a canal supplying water for three different crops . Maximum capacity of the canal is Q units of water. Amount of water allocated to each field as  $x_i$ .



Net benefits (NB) from producing the crops can be expressed as a function of the water allotted.

$$NB_1(x_1) = 5x_1 - 0.5x_1^2$$

$$NB_2(x_2) = 8x_2 - 1.5x_2^2$$

$$NB_3(x_3) = 7x_3 - x_3^2$$

### Optimization Problem:

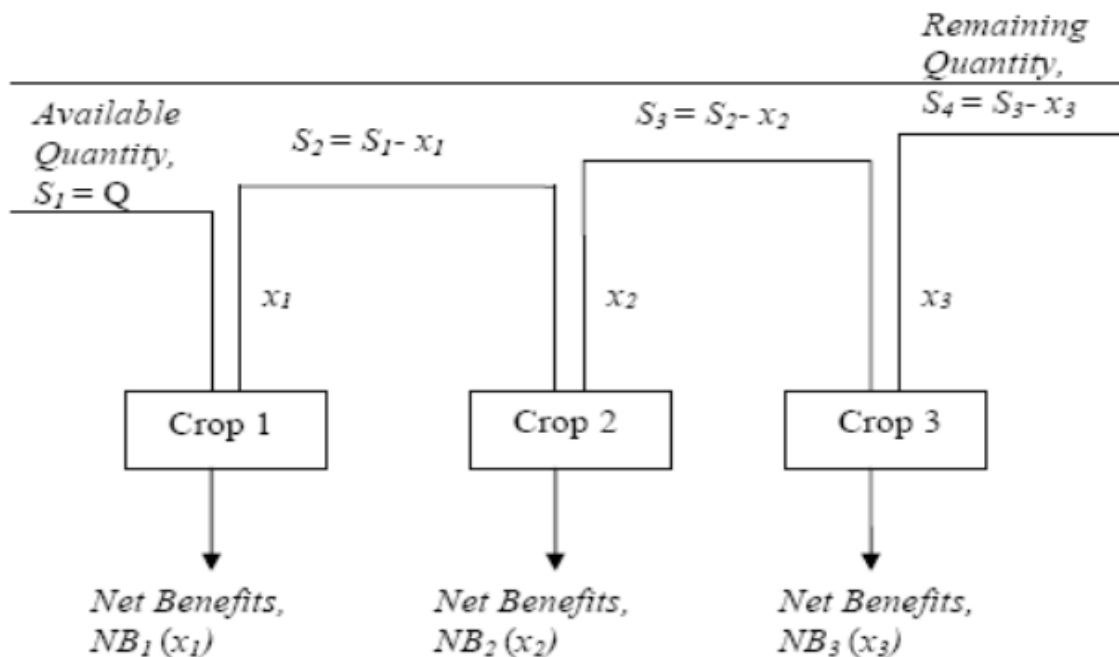
Determine the optimal allocations  $x_i$  to each crop that maximizes the total net benefits from all the three crops

Structure the problem as a sequential allocation process or a multistage decision making procedure. Allocation to each crop is considered as a decision stage in a sequence of decisions. Amount of water allocated to crop  $i$  is  $x_i$ . Net benefit from this allocation is  $NB_i(x_i)$ . Introduce one state variable  $S_i$  :- Amount of water available to the remaining  $(3-i)$  crops. State transformation equation can be written as

$$S_{i+1} = S_i - x_i$$

### Sequential Allocation Process

The allocation problem is shown as a sequential process



Backward Recursive Equations Objective function:-- To maximize the net benefits

$$\max \sum_{i=1}^3 NB_i(x_i)$$

Subjected to the constraints

$$\begin{aligned} x_1 + x_2 + x_3 &\leq Q \\ 0 \leq x_i &\leq Q \quad \text{for } i = 1, 2, 3 \end{aligned}$$

Let  $f_1(Q)$  be the maximum net benefits that can be obtained from allocating water to crops 1, 2 and 3

$$f_1(Q) = \max_{\substack{x_1 + x_2 + x_3 \leq Q \\ x_1, x_2, x_3 \geq 0}} \left[ \sum_{i=1}^3 NB_i(x_i) \right]$$

Transforming this into three problems each having only one decision variable

$$f_1(Q) = \max_{\substack{x_1 \\ 0 \leq x_1 \leq Q}} \left[ NB_1(x_1) + \max_{\substack{x_2 \\ 0 \leq x_2 \leq Q - x_1 = S_2}} \left\{ NB_2(x_2) + \max_{\substack{x_3 \\ 0 \leq x_3 \leq S_2 - x_2 = S_3}} NB_3(x_3) \right\} \right]$$

Now starting from the last stage, let  $f_3(S_3)$  be the maximum net benefits from crop 3.

State variable  $S_3$  for this stage can vary from 0 to Q

Thus,

$$f_3(S_3) = \max_{\substack{x_3 \\ 0 \leq x_3 \leq S_3}} NB_3(x_3)$$

But  $S_3 = S_2 - x_2$ . Therefore  $f_3(S_3) = f_3(S_2 - x_2)$

Hence,

$$f_1(Q) = \max_{0 \leq x_1 \leq Q} \left[ NB_1(x_1) + \max_{0 \leq x_2 \leq Q - x_1 = S_2} \{NB_2(x_2) + f_3(S_2 - x_2)\} \right]$$

Now, let  $f_2(S_2)$  be the maximum benefits derived from crops 2 and 3 for a given quantity  $S_2$  which can vary between 0 and Q

Therefore,

$$f_2(S_2) = \max_{0 \leq x_2 \leq Q - x_1 = S_2} \{NB_2(x_2) + f_3(S_2 - x_2)\}$$

Now since  $S_2 = Q - x_1$ ,  $f_1(Q)$  can be rewritten as

$$f_1(Q) = \max_{0 \leq x_1 \leq Q} [NB_1(x_1) + f_2(Q - x_1)]$$

Once the value of  $f_3(S_3)$  is calculated, the value of  $f_2(S_2)$  can be determined, from which  $f_1(Q)$  can be determined.

Let the function  $f_i(S_i)$  be the total net benefit from crops 1 to  $i$  for a given input of  $S_i$  which is allocated to those crops.

Considering the first stage,

$$f_1(S_1) = \max_{0 \leq x_1 \leq S_1} NB_1(x_1)$$

Solve this equation for a range of  $S_1$  values from 0 to Q

Considering the first two crops, for an available quantity of  $S_2$ ,  $f_2(S_2)$  can be written as

$$f_2(S_2) = \max_{0 \leq x_2 \leq S_2} [NB_2(x_2) + f_1(S_2 - x_2)]$$

$S_2$  ranges from 0 to Q

Considering the whole system,  $f_3(S_3)$  can be expressed as,

$$f_3(S_3) = \max_{x_3 \leq S_3 = Q} [NB_3(x_3) + f_2(S_3 - x_3)]$$

If the whole Q units of water should be allocated then the value of  $S_3$  can be taken as equal to Q

Otherwise,  $S_3$  will take a range of values from 0 to Q

The basic equations for the water allocation problem using both the approaches are discussed

Consider the example previously discussed with the maximum capacity of the canal as **4 units**. The net benefits from producing the crops for each field are given by the functions below.

$$NB_1(x_1) = 5x_1 - 0.5x_1^2$$

$$NB_2(x_2) = 8x_2 - 1.5x_2^2$$

$$NB_3(x_3) = 7x_3 - x_3^2$$

The possible net benefits from each crop are calculated according to the functions given and are given in Table 1.

**Table 1**

$x_i$	$NB_1(x_1)$	$NB_2(x_2)$	$NB_3(x_3)$
0	0.0	0.0	0.0
1	4.5	6.5	6.0
2	8.0	10.0	10.0
3	10.5	10.5	12.0
4	12.0	8.0	12.0

$$f_3(S_3) = \max NB_3(x_3) \text{ with the range of } S_3 \text{ from 0 to 4.}$$

The calculations for this stage are shown in the table 2.

**Table 2**

State $S_3$	$NB_3(x_3)$						$f_3(S_3)$	$x_3^*$
	$x_3:$	0	1	2	3	4		
0		0					0	0
1		0	6				6	1
2		0	6	10			10	2
3		0	6	10	12		12	3
4		0	6	10	12	12	12	3,4

Next, by considering last two stages together, the suboptimization function is

$$f_2(S_2) = \max [NB_2(x_2) + f_1(S_2 - x_2)]$$

This is solved for a range of  $S_2$  values from 0 to 4. The value of  $f_3(S_2 - x_2)$  is noted from the previous table. The calculations are shown in Table 3.

**Table 3**

State $S_2$	$x_2$	$NB_2(x_2)$	$(S_2 - x_2)$	$f_3(S_2 - x_2)$	$f_2(S_2) =$ $NB_2(x_2) +$ $f_3(S_2 - x_2)$	$f_2^*(S_2)$	$x_2^*$
0	0	0	0	0	0	0	0
1	0	0	1	6	6	6.5	1
	1	6.5	0	0	6.5		
2	0	0	2	10	10	12.5	1
	1	6.5	1	6	12.5		
	2	10	0	0	10		
3	0	0	3	12	12	16.5	1
	1	6.5	2	10	16.5		
	2	10	1	6	16		
	3	10.5	0	0	10.5		
4	0	0	4	12	12	20	2
	1	6.5	3	12	18.5		
	2	10	2	10	20		
	3	10.5	1	6	16.5		
	4	8	0	0	8		

Finally, by considering all the three stages together, the sub-optimization function is

$$f_1(Q) = \max_{0 \leq x_1 \leq Q} [NB_1(x_1) + f_2(Q - x_1)].$$

The value of  $S_1 = Q = 4$ . The calculations are shown in

the table 4.

**Table 4**

State	$x_1$	$NB_1(x_1)$	$(Q-x_1)$	$f_2(Q-x_1)$	$f_1(S_1)=$ $NB_1(x_1)+$ $f_2(Q-x_1)$	$f_1^*(S_1)$	$x_1^*$
$S_1 = Q$	0	0	4	20	20		
	1	4.5	3	16.5	21		
	2	8	2	12.5	20.5	21	1
	3	10.5	1	6.5	17		
	4	12	0	0	12		

Now, backtracking through each table to find the optimal values of decision variables, the optimal allocation for crop 1,  $x_1^* = 1$  for a  $S_1$  value of 4. This will give the value of  $S_2$  as  $S_2 = S_1 - x_1 = 3$ . From Table 3, the optimal allocation for crop 2,  $x_2$  for  $S_2 = 3$  is 1. Again,  $S_3 = S_2 - x_2 = 2$ . Thus,  $x_3^*$  from Table 2 is 2. The maximum total net benefit from all the crops is 21. The optimal solution is given below.

$$f^* = 21$$

$$x_1^* = 1$$

$$x_2^* = 1$$

$$x_3^* = 2$$

$$f_1(Q) = \max_x [NB_1(x_1) + f_2(Q - x_1)] \text{ -----1}$$

$$f_2(S_2) = \max_x [NB_2(x_2) + f_1(S_2 - x_2)] \text{ -----2}$$

$$f_3(S_3) = \max NB_3(x_3) \text{ -----3}$$

Equations (1), (2) and (3) are called as the recursive set of Equations.



### **Tutorial Class problem**

(Q) Consider that a quantity of water= $Q$  has to be allotted to three users denoted by  $j=1,2,3$  where ( $j$ ) is the user. The problem is to supply three water quantities  $X_1, X_2$  and  $X_3$  to three users 1,2 and 3 so as to maximize the total net benefits.